Review Material, Exam 1, Ops Research

The exam will cover material from Chapter 3 (we skipped 3.6, 3.7), up through section 4.8. Section 4.9 may be used to help with LINDO, and we skipped 4.10, but we did look at 4.11 (degeneracy).

Background Material: Linear Algebra

Though these may not be asked explicitly, you should be able to do the following (and may be part of solutions):

- 1. Be able to solve $A\mathbf{x} = \mathbf{b}$ for different types of matrices (if A is square and invertible, and if it is tall and full rank, tall and not full rank, wide and full rank, wide and not full rank).
- 2. Be able to discuss the solution to $A\mathbf{x} = \mathbf{b}$ in terms of the row space, the column space and the null space of A.
- 3. Be able to compute the determinant and inverse of a matrix A
- 4. Note that our book uses "ERO" for elementary row operation.
- 5. Define what it means for a set of vectors to be linearly independent.
- 6. Be able to solve for a basis for the row space, column space and null space for a given matrix A.

Definitions:

Linear program, standard form of a linear program, objective function, constraints (binding and non-binding), extreme point, isoprofit line, feasible region, unrestricted variable (URS), BFS, adjacent BFS, slack/surplus (or excess) variables, basic solution, BFS, BV, NBV, direction of unboundedness, Convex set, convex combination.

The word "degenerate" is used in two distinct cases: See the top of p. 134 for what it means for a linear program to be degenerate. In class, we said that a BFS is degenerate if one of the basic variables (for the optimal solution) is zero.

Skills

- 1. Translate unrestricted variables so that all variables are non-negative.
- 2. Be able to show that a given set is convex, use convexity in other arguments (See review questions below for examples).
- 3. Be able to set up and solve an LP graphically. Given the objective function, be able to state the direction in which the function increases (or decreases) the fastest.
- 4. Be able to set up an LP generally, and of specific types: A diet problem, a work scheduling problem, a production process model, blending, and multiperiod problems (like the sailboat example).
- 5. Be able to translate a LP into standard form given "less than or equal to" constraints, "greater than or equal to" constraints, and change the variables so that they are all non-negative.
- 6. Be able to write the simplex tableau from the linear program (for both a maximization and a minimization problem).
- 7. Be able to solve a linear program (a maximization problem with "less than or equal to" constraints) using the Simplex Method.
- 8. Given a tableau, be able to tell if it is a terminal tableau, and interpret what the solution is (unique, multiple, unbounded).
- 9. (For something that might be take home) Be able to use either LINGO or a spreadsheet program to solve an LP.

Theorems:

- Theorem 1 says that extreme points are the same BFS (and are the same as corner points).
- Theorem 2: The Representation Theorem Says that any feasible point can be written as

$$\mathbf{d} + \sum_{i=1}^{N} \sigma_i \mathbf{b}_i$$

where \mathbf{d} is a direction of unboundedness and the sum is a convex combination of the corner points of the feasible set.

• Theorem 3: The Fundamental Theorem of Linear Programming

If an LP has an optimal solution, then it has an optimal BFS (or corner point, or extreme point).

Review Questions

- 1. What are the four possible outcomes when solving a linear program? Hint: The first is that there is a unique solution to the LP.
- 2. The following are to be sure you understand the process of constructing a linear program:
 - (a) Draw a production process diagram and set up the LP for Exercise 6, p. 98 (Sect. 3.9).
 - (b) Exercise 2, 31 Chapter 3 review (Be sure you can solve an LP graphically)
 - (c) Exercise 6, 18 Chapter 3 review (A ton is 2000 lbs)
 - (d) Exercise 22, Chapter 3 review. Hint: Consider using a triple index on your variables.
 - (e) Exercise 47, 53 in Chapter 3 review.
- 3. Convert the following LP to one in standard form. Write the result in matrix-vector form, giving \mathbf{x} , \mathbf{c} , A, \mathbf{b} (from our formulation).

$$\min z = 3x - 4y + 2z$$

$$\text{st} \quad 2x - 4y \ge 4$$

$$x + z \ge -5$$

$$y + z \le 1$$

$$x + y + z = 3$$

with $x \ge 0$, y is URS, $z \ge 0$.

4. Consider again the "Wyndoor" company example we looked at in class:

$$\min z = 3x_1 + 5x_2$$
st $x_1 \le 4$
 $2x_2 \le 12$
 $3x_1 + 2x_2 \le 18$

with x_1, x_2 both non-negative.

- (a) Rewrite so that it is in standard form.
- (b) Let s_1, s_2, s_3 be the extra variables introduced in the last answer. Is the following a basic solution? Is it a basic feasible solution?

$$x_1 = 0, x_2 = 6, s_1 = 4, s_2 = 0, s_3 = 6$$

Which variables are BV, and which are NBV?

- (c) Find the basic feasible solution obtained by taking s_1, s_3 as the non-basic variables.
- 5. Consider Figure 1, with points A(1,1), B(1,4) and C(6,3), D(4,2) and E(4,3).

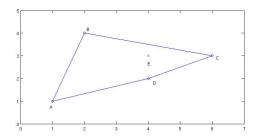


Figure 1: Figure for the convex combinations, Exercise 6.

- Write the point E as a convex combination of points A, B and C.
- Can E be written as a convex combination of A, B and D? If so, construct it.
- Can A be written as a linear combination of A, B and D? If so, construct it.
- 6. Draw the feasible set corresponding to the following inequalities:

$$x_1 + x_2 \le 6$$
, $x_1 - x_2 \le 2$ $x_1 \le 3$, $x_2 \le 6$

with x_1, x_2 non-negative.

- (a) Find the set of extreme points.
- (b) Write the vector $[1,1]^T$ as a convex combination of the extreme points.
- (c) True or False: The extreme points of the region can be found by making exactly two of the constraints binding, then solve.
- (d) If the objective function is to maximize $2x_1 + x_2$, then (a) how might I change that into a minimization problem, and (b) solve it.
- 7. Consider the unbounded feasible region defined by

$$x_1 - 2x_2 \le 4$$
, $-x_1 + x_2 \le 3$

with x_1, x_2 non-negative. Consider the vector $\mathbf{p} = [5, 2]$.

- (a) Show that **p** is in the feasible region.
- (b) Set up the system you would solve in order to write **p** in the form given in Theorem 2 above (provide a specific vector **d**).
- 8. Finish the definition: Two basic feasible solutions are said to be adjacent if:
- 9. Let **d** be a direction of unboundedness. Using the *definition*, prove that this means that $r\mathbf{d}$ is also a direction of unboundedness, for any constant $r \ge 0$.
- 10. If C is a convex set, then $\mathbf{d} \neq 0$ is a direction of unboundedness for C iff $\mathbf{x} + d \in C$ for all $\mathbf{x} \in C$ (Use the definition of unboundedness).
- 11. For an LP in standard form (see above), prove that the vector \mathbf{d} is a direction of unboundedness iff $A\mathbf{d} = 0$ and $\mathbf{d} > 0$.
- 12. Show that the set of optimal solutions to an LP (assume in standard form) is convex.
- 13. Let a feasible region be defined by the system of inequalities below:

$$\begin{array}{rcl}
-x_1 + 2x_2 & \leq 6 \\
-x_1 + x_2 & \leq 2 \\
x_2 & \geq 1 \\
x_1, x_2 > 0
\end{array}$$

The point (4,3) is in the feasible region. Find vectors \mathbf{d} and $\mathbf{b}_1, \dots \mathbf{b}_k$ and constants σ_i so that the Representation Theorem is satisfied (NOTE: Your vector \mathbf{x} from that theorem is more than two dimensional).

14. Let a feasible region be defined by the system of inequalities below:

$$\begin{array}{rcl}
-x_1 + x_2 & \leq 2 \\
x_1 - x_2 & \leq 1 \\
x_1 + x_2 & \leq 5 \\
x_1, x_2 \geq 0
\end{array}$$

The point (2,2) is in the feasible region. Find vectors \mathbf{d} and $\mathbf{b}_1, \dots \mathbf{b}_k$ and constants σ_i so that the Representation Theorem is satisfied (NOTE: Your vector \mathbf{x} from that theorem is more than two dimensional).

15. Suppose that $\lambda_1 \leq \lambda_2 \leq \cdots \lambda_n$, and let $\sigma_1, \cdots, \sigma_n$ be non-negative constants so that $\sum_{i=1}^n \sigma_i = 1$. Show that

$$\lambda_1 \le \sigma_1 \lambda_1 + \sigma_2 \lambda_2 + \dots + \sigma_n \lambda_n \le \lambda_n$$

16. Show that, if **x** is in the convex hull of vectors $\mathbf{b}_1, \cdots \mathbf{b}_k$, then for any constant vector **c**,

$$\mathbf{c}^T \mathbf{x} \leq \max_i \left\{ \mathbf{c}^T \mathbf{b}_i \right\}$$

- 17. True or False, and explain: The Simplex Method will always choose a basic feasible solution that is adjacent to the current BFS.
- 18. Given the current tableau (with variables labeled above the respective columns), answer the questions below.

- (a) Is the tableau optimal (and did your answer depend on whether we are maximizing or minimizing)? For the remaining questions, you may assume we are maximizing.
- (b) Give the current BFS.
- (c) Directly from the tableau, can I increase x_2 from 0 to 1 and remain feasible? Can I increase it to 4?
- (d) If x_2 is increased from 0 to 1, compute the new value of z, x_1, s_1 (assuming s_2 stays zero).
- (e) Write the objective function and all variables in terms of the non-basic (or free) variables, and then put them in vector form.