

Chapter 7 Review Solutions

Exercises 1-6, 8-11, 15-16, 23, 25

- After using the NW corner rule and apply MODI (modified distribution), or the “u-v” method, we get:

	C1	C2	C3	
P 1	20	90	30	50
P 2	10	90	0	100
P 3	50	0	0	50
Dummy	0	0	70	70
Demand	80	90	100	

- You should find one optimal way of performing the assignments is:
Worker 1 does job 3. Worker 2 does job 2. Worker 3 does nothing. Worker 4 does job 1. Worker 5 does job 4.
- To set this up, think of Jan, Feb and March having both a supply and demand- The supplies being 35, 30 and 35, respectively, and the demands being 30, 30, 20, 20, respectively.

Here is the optimal solution:

	Jan	Feb	Mar	Dummy	
Jan	30	5	0	0	35
Feb	0	10	0	20	30
Mar	0	15	20	0	35
Demand	30	30	20	20	

4. For the maid assignment problem, we should find that we may assign:
Maid 1 to the bathroom, Maid 2 to straighten up, Maid 3 to the kitchen, Maid 4 is idle, and Maid 5 to vacuum.
5. To set it up, think of the hard drive (HD), the computer memory (CM), and the computer tape (CT) as being the supplies- 200, 100 and 300 each. Then they (computer packets) can be “delivered” as word processing (WP), packaged program (PP), data file (DF), with “demands” 300, 100 and 100. The costs are costs in time. Notice that Table 68 does need to be modified for our final costs- For example, if WP is accessed 8 times a month, then the time costs for the HD, CM and CT are $8 \times [5, 2, 10]$, or $[40, 16, 80]$. Here is the optimal tableau:

	WP	CM	CT	Dummy	
WP	40	16	8	0	200
	200	0			
PP	16	4	2	0	100
	100				
DF	80	32	12	0	300
		100	100	100	
Demand	300	100	100	100	

6. In setting this problem up, we’ll need two “dummy” calls. Here is one ending array from which we could get an assignment (cars down the side, calls across the rows)

4	6	13	0	0	→	Car 1 to Dummy 1
0	2	2	0	0		Car 2 to Dummy 2
1	3	0	0	0		Car 3 to Call 3
0	2	0	1	1		Car 4 to Call 1
4	0	3	1	1		Car 5 to Call 2

8. Notice that we needed a dummy column to proceed. An optimal solution: $x_{11} = 40$, $x_{12} = 10$, $x_{22} = 50$, $x_{32} = 10$, and $x_{34} = 30$.
9. The LP is equivalent to a 2×2 transportation problem. Taking x_1, x_2 for the first row and x_3, x_4 for the second row, the optimal solution is 3, 1, 0, 5, respectively.
10. You should find that the optimal solution is given by

10	5		$\Rightarrow z = 130$
	5	10	

11. (TYPO: s_i should be $s_1 = 16$). Here's the new optimum:

	7	2	4	16
10	70	6		
	12	8	4	15
		4	11	
	10	10	11	

$Z = 40 + 12 + 32 + 44 = 128$ ← Shipping 1 unit more cost - 2 dollars less.

1 unit switched from X_{22} to X_{12} (savings: \$6)
 1 unit more went via X_{23} (cost: \$4) } saving \$2

15. Since the change is to the cost for a NBV, changing it will only result in changing that value of $c_{ij} - (u_i + v_j)$, which in this case:

$$(7 + \Delta) - (3 + 2) = 2 + \Delta > 0 \Rightarrow \Delta > -2$$

Therefore, the value of c_{24} can range in the interval $[5, \infty)$

NOTE: It's OK to leave the value of 0 out- This would permit a multiple optimal solution, but the current solution would still be one of them.

16. This changes the cost associated to a basic variable, so we will have to re-work the our u, v computations. Solving the resulting inequalities should yield $-2 \leq \Delta \leq 2$, so the current basis remains optimal for $11 \leq c_{23} \leq 15$.

$v_1 = 6 - \Delta \quad v_2 = 6 \quad v_3 = 10 \quad v_4 = 2$

	8	6	10	9
$u_1 = 0$	$(2 + \Delta)$	10	25	(7)
$u_2 = 3 + \Delta$	9	12	$13 + \Delta$	7
	14	9	16	5
$u_3 = 3$	$(5 + \Delta)$	10	(3)	30

23. Two notes that might help you construct this one- Since customer three has two demands, split it into two columns- one for the regular demand, and one for the extra demand. We will need a dummy plant to make the problem balanced.

Further, we are going to maximize the profit in this problem, not minimize. Therefore, we don't want to use M for the "profit" from the dummy plant to the customer- Rather, we can use $-M$. The tableau is given below- We have given supply and demand in thousands.

	C_1	C_2	C_{31}	C_{32}	C_4	
P_1	65	63	62	62	64	3
P_2	68	67	65	65	62	5
P_3	63	60	59	59	60	5
D	$-M$	$-M$	$-M$	0	0	3
	4	3	3	3	3	

25. In setting this up, customer 1 gets two rows and we'll need a dummy column (with zero costs). The second problem is that this is a maximization problem, not a minimization problem- Therefore, make the bids in Table 84 negative values, then we can minimize.

This is only different until you "subtract" the row minimums, at which point the array becomes non-negative again (and we proceed normally- See below). At the optimal value, Customer 1 should buy painting 1 (and the dummy), Customer 2 should buy Painting 2, Customer 3 should buy painting 3, and customer 4 gets painting 4 (the diagonal elements!)

The image shows a handwritten tableau illustrating the process of subtracting row minimums from a cost matrix to make it non-negative. The initial cost matrix is shown on the left, with row minimums listed to its right. An arrow points to the resulting matrix on the right, where the row minimums have been subtracted from each row. The resulting matrix is non-negative, and the row minimums are listed below it.

C_1	C_1'	C_2	C_3	C_4
-8	-11	0	0	0
-8	-11	0	0	0
-9	-13	-12	-7	0
-9	0	-11	0	0
0	0	-12	-9	0

Row mins:
-11
-11
-13
-11
-12

→

3	0	11	11	11
3	0	11	11	11
4	0	1	6	13
2	11	0	11	11
12	12	0	3	12
2	0	0	3	11

etc.

← Col min