

# What Makes Writing in Mathematics “Good”?

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# Intrinsic Dimensionality Estimation via Singular Value Curves

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## Abstract

We propose a modification and a generalization of a method for determining the intrinsic dimensionality (ID) of data, the method of singular value (SV) curves. While other methods use a single computation of the covariance of a ball of data, this method takes multiple measurements of a ball of increasing radius, and is therefore itself a generalization of the widely used PCA spectrum. Our modification is geared towards making the SV curves canonical, allowing for an automation of the ID estimation. We then generalize the technique to give an analytic version of the SV curves, making them suitable to analyze noisy data sets. We give several examples and show that the extra computational effort is well worthwhile when one desires an in-depth analysis of the local topological structure of the data.

## 1 Introduction

Local representations of smooth data are linear. In fact, they result directly from the well-known Taylor's formula in  $m$ -dimensions which provides a basis for the tangent space about the point of approximation. This tangent space is the column space of the Jacobian of the surface function which, in data analysis problems, is generally unknown. Here we want to be explicit about our assumptions on the data. We will assume that a given (local) data set, has the following decomposition,

$$x \in \mathbb{R}^n \Rightarrow \tilde{x} \in \mathbb{R}^k$$

where there exists a (perhaps nonlinear) function  $f$  so that  $x = f(\tilde{x})$  and  $k$  is the smallest such integer. Thus, our  $\tilde{x}$  is a  $k$ -dimensional parametrization of the surface given as data points  $x$ , and  $k$  is the intrinsic dimension of the data. Of course, in practice, the equality given above is only an approximation.

We note that this is in contrast to other definitions of intrinsic dimension, notably from [1], where the definition is given as the number of dimensions necessary to *encapsulate* the data- if the intrinsic dimension of the data is  $k$  by this definition, then the data lies entirely

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- Notation consistent and readable.

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- “Useful and interesting” does not tell the reader *why* or *where*.

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Let  $G$  be a finite graph. A *dominating set* for  $G$  is a subset  $D$  of vertices such that every vertex of  $G$  is either in  $D$  or adjacent to a vertex in  $D$ .

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**Proposition.** If  $n$  is even, then  $n^2$  is divisible by 4.

**Proof.**  $n$  is even so  $n = 2k$ . Then  $n^2 = 4k^2$  which is divisible by 4. Done.

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# Quick Checklist for Student Drafts

When revising, ask:

- 1 Did I define every nonstandard term and symbol *before* using it?
- 2 Did I tell the reader what the goal is (theorem, classification, method, example set)?
- 3 Does each proof have a visible structure (strategy + justified steps)?
- 4 Do my equations have “glue sentences” that explain what they mean?
- 5 Could a peer summarize each paragraph’s purpose in one sentence?



# Wrap Up

- In pairs, rewrite the "bad" examples into "good" examples.
- Review each other's work
- Rewrite.
- Turn in your papers at the end of class.